

# Reasoning about Trust and Aboutness in the Context of Communication

Robert Demolombe

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- ▶ The truster  $i$  trusts the trustee  $j$  in some property about proposition  $p$
- ▶ Giulietta trusts Romeo in his sincerity about the fact that Giulietta is very pretty
- ▶ Giulietta believes that:  
IF Romeo told her that she is very pretty,  
THEN Romeo believes that she is very pretty

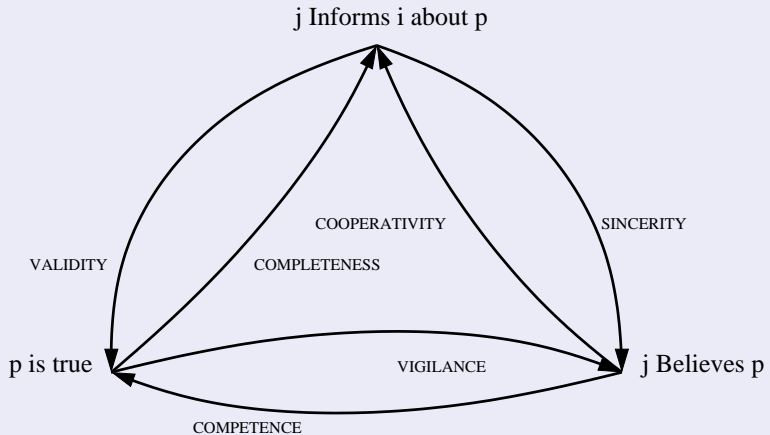
## Trusted properties

Form: IF ... Then ...

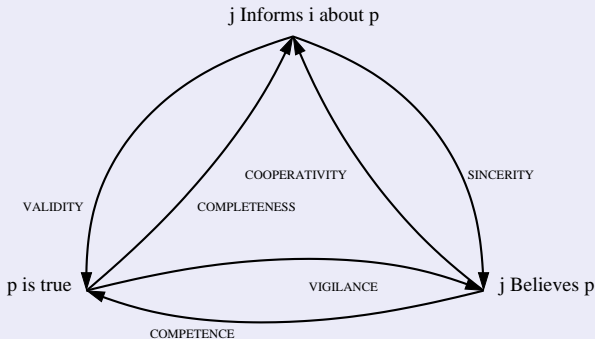
Relationships between:

- ▶  $j$  Informs  $i$  about  $p$  :  $Inf_{j,i}(p)$
- ▶  $j$  Believes  $p$  :  $Bel_j(p)$
- ▶  $p$  is true :  $p$

## Trusted properties



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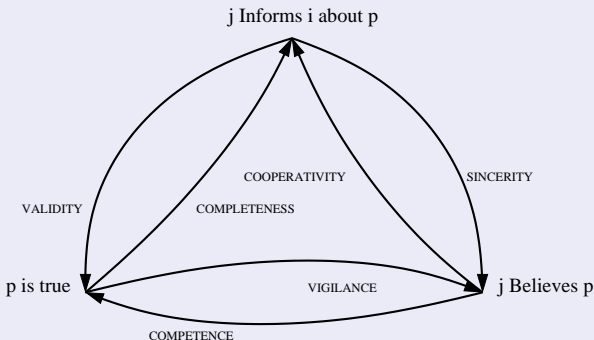
$p$ : Giulietta is very pretty

SINCERITY: If Romeo Informs Giulietta about  $p$ , Then Romeo Believes  $p$

COMPETENCE: If Romeo Believes  $p$ , Then  $p$  is true

VALIDITY: If Romeo Informs Giulietta about  $p$ , Then  $p$  is true

## Trusted properties



$p$ : Giulietta is very pretty

VIGILANCE: If  $p$  is true, Then Romeo Believes  $p$

COOPERATIVITY: If Romeo Believes  $p$ , Then Romeo Informs Giulietta about  $p$

COMPLETENESS: If  $p$  is true, Then Romeo Informs Giulietta about  $p$



## Formalization

$p$ : Classical Propositional Calculus

$Bel_i(p)$ : KD logic

$Inf_{j,i}(p)$ : Classical Modal (not Normal) Logic

$TrustSinc(i, j, p) \stackrel{\text{def}}{=} Bel_i(Inf_{j,i}(p) \rightarrow Bel_j(p))$

Assumption: perfect communication

(OBS)  $Inf_{j,i}\phi \rightarrow Bel_i(Inf_{j,i}\phi)$

(OBS')  $\neg Inf_{j,i}\phi \rightarrow Bel_i(\neg Inf_{j,i}\phi)$

## Formal example

$q$  = Julietta is in Toulouse

Romeo trusts Julietta in her completeness about  $q$

$Bel_R(q \rightarrow Inf_{G,R}(q))$

$\neg Inf_{G,R}(q)$

Entails:  $Bel_R(\neg q)$

## Formal example

$q =$  Julietta is in Toulouse

Romeo trusts Julietta in her completeness about  $q$

$Bel_R(q \rightarrow Inf_{G,R}(q))$

$\neg Inf_{G,R}(q)$

Entails:  $Bel_R(\neg q)$

Romeo trusts X in his sincerity about  $q$

$Bel_R(Inf_{X,R}(q) \rightarrow Bel_X(q))$

$Inf_{X,R}(q)$

Entails:  $Bel_R(Bel_X(q))$

Entails:  $Bel_R(\neg q \wedge Bel_X(q))$

## Trust extended to topics (with Andrew J.I. Jones)

- ▶ Giulietta trusts Romeo in his validity about every sentence about the topic modal logic
- ▶ If sentence ' $p$ ' is about modal logic, then Giulietta trusts Romeo in his validity about the truth of the proposition represented by  $p$

What does it mean that sentence ' $p$ ' is about topic  $t$ ?

## Trust extended to topics

"sentence ' $p$ ' is about topic  $t$ "

does not depend on the truth value of proposition  $p$

it depends on its meaning

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- ▶ "KD is decidable" is about modal logic
- ▶ "KD is not decidable" is about modal logic
- ▶ "KD is decidable or KD is not decidable" is about modal logic
- ▶ "Toulouse is in Spain or Toulouse is not in Spain" is NOT about modal logic

## Formalization

$A(t, p')$ : the proposition named by  $p'$  is about the topic  $t$   
Language of CPC.

atoms of CPC +  $A(t, p')$

Models. Two sorts for the interpretation of  $A(t, p')$

Satisfiability. Bochvar's 3 valued logic: true, false, undefined

$\text{true} \vee \text{undefined} = \text{undefined}$

## Formalization

Valid schema:

If  $\models p \leftrightarrow q$  and ' $p'$ ' and ' $q'$ ' same atoms

Then  $\models A(t, p') \leftrightarrow A(t, q')$

Additional schema:

$A(t, p') \rightarrow A(t, \neg p')$



## Formalization

Valid schema:

If  $\models p \leftrightarrow q$  and  $'p'$  and  $'q'$  same atoms

Then  $\models A(t, 'p') \leftrightarrow A(t, 'q')$

Additional schema:

$A(t, 'p') \rightarrow A(t, '\neg p')$

Rejected schema:

$A(t, 'p \wedge q') \rightarrow A(t, 'p') \vee A(t, 'q')$

Example:  $t$ : bigamy

$'p'$ : Romeo is married with Giulietta

$'q'$ : Romeo is married with Venus

## Trust extended to individuals (with Luis Fariñas del Cerro)

- ▶ Romeo trusts Giulietta in her validity about all the sentences which inform about Sorolla
- ▶ If  $p$  informs about Sorolla, then Romeo trusts Giulietta in her validity about  $p$

What does it mean that  $p$  informs about Sorolla?

## Intuition

$p$  = Sorolla is a Spanish painter

$p$  informs about Sorolla

$\neg p$  informs about Sorolla

$p \vee \neg p$  does not inform about Sorolla

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$p$  = Sorolla is a Spanish painter

$p$  informs about Sorolla

$\neg p$  informs about Sorolla

$p \vee \neg p$  does not inform about Sorolla

$q$  = Sorolla is a Spanish painter or Picasso is a Spanish painter

$q$  informs about Sorolla

## Formalization

$p$  does NOT inform about Sorolla iff

In any model

If we change the truth value of the tuples which contain  $d$  and  $d$  is the interpretation of Sorolla (or a term which contains Sorolla),

Then the truth value of  $p$  does not change

(except if  $d$  is also the interpretation of another term)

## Topics and Individuals

Giulietta trusts Romeo in his validity about the sentences which *inform about Gödel* and which are *about the topic* Logic

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### **Possible extension**

Sentences which inform about a relationship between two individuals

Romeo loves Giulietta: YES

Romeo leaves in Toulouse and AND Giulietta leaves in Chiclana:  
NO

There exist a city  $x$  (Romeo leaves in  $x$  AND Giulietta leaves in  $x$ ):  
YES

## Graded Trust (with Leyla Amgoud)

What does it mean that  $i$  strongly trusts  $j$ ?

Strength of  $i$ 's belief?

Regularity level of  $Inf_{j,i}(p) \rightarrow Bel_j(p)$ ?

Combination of both:

$Bel_i^g(Inf_{j,i}(p) \Rightarrow^h Bel_j(p))$



## Conclusion

Many interesting topics to be investigated

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